

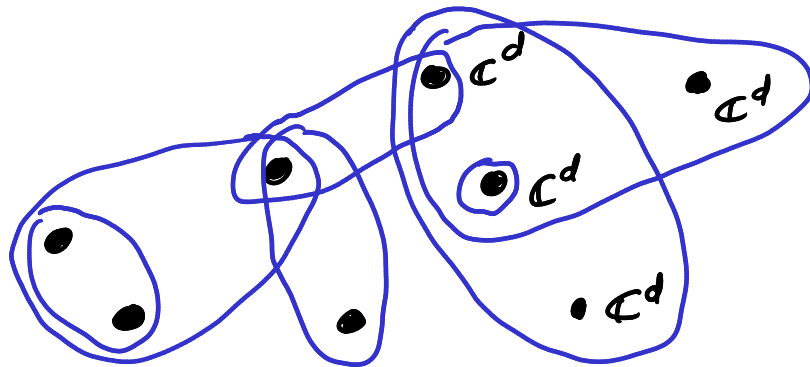
# I. Hamiltonian Complexity

## Notation & Terminology

- $\lambda_0(H)$  or  $\lambda_{\min}(H) := \min_{|\phi\rangle} \langle \phi | H | \phi \rangle$   
"ground state energy"
- $H |\psi_0\rangle = \lambda_0 |\psi_0\rangle$      $|\psi_0\rangle =$  "ground state"
- $\mathcal{L}_0 = \text{span} \{ |\psi\rangle : H|\psi\rangle = \lambda_{\min} |\psi\rangle \}$   
 $\mathcal{L}_0 =$  "g.s. subspace"
- Note:  $H \geq 0$ ,  $\lambda_{\min}(H) = 0 \Rightarrow$  g.s. subspace =  $\ker H$
- $\Delta(H) = \lambda_1 - \lambda_0 = \min_{|\phi\rangle \perp \mathcal{L}_0} \langle \phi | H | \phi \rangle - \lambda_0(H)$   
"spectral gap": same operator, different eigvals
- $\lambda_0(H_1) - \lambda_0(H_2)$   
"promise gap": different operators, same eigvals

# Local Hamiltonians

Core object of study for this entire course.



Many-body quantum system:

- Multipartite Hilbert space  $\mathcal{H} = (\mathbb{C}^d)^{\otimes n}$   
(often  $d=2$ :  $n$  qubit system)  
e.g. particles with spin- $s$  ( $d = \frac{s(s+1)}{2}$ )

- Local interaction Hamiltonian:

$$h = h_S \otimes \mathbb{1}_{[n] \setminus S}, \quad S \subset [n]$$

i.e.  $h$  acts non-trivially only on subset  $S$  of the particles

We say that interaction  $h$  is  $k$ -local if  $|S| = k$  (i.e. acts non-trivially on  $k$  particles).

Notation: often write " $h_S$ "  $\equiv h_S \otimes \mathbb{1}_{[n] \setminus S}$ , i.e. subscript indicates indices of particles acted on,  $\mathbb{1}$  on rest is implicit.


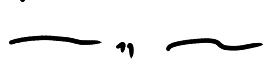
Def ( $k$ -local Hamiltonian)

$$H = \sum_i h^{(i)} \in \mathcal{B}(\mathbb{C}^d)^{\otimes n}$$

We say that  $H$  is a  $k$ -local Hamiltonian (or " $H$  is  $k$ -local") if  $\forall i$   $h^{(i)}$  is  $k$ -local.

I. e.  $k$ -local many-body Hamiltonian made of many interactions, each involving at most  $k$  particles.

Example: spins on line with nearest-neighbour interactions:  $H = \sum_{i=1}^{n-1} h_{i,i+1}$

spins on lattice   
 :  $H = \sum_{\langle i,j \rangle} h_{ij}$   
↑ neighbours

Note: In general, **no** underlying geometry  $\Rightarrow$  **no** requirement that local interactions are geometrically local.

Physics terminology: " $k$ -local" = " $k$ -body" but stuck with " $k$ -local" now in QIT.

Motivation:

Fundamental interactions in nature (e.g. electromagnetic, or "Coulomb", interaction) are not  $k$ -local — interactions decay as  $\frac{1}{\text{poly}}$ , not finite-range.

However, in many-body system electron clouds of neighbouring atoms shield atoms from fields of more distant particles (cf. Faraday cage).

At low-energies, many-body systems often well-approximated by  $k$ -local (often 2-local) Hamiltonians.

→  $k$ -local  $H$  crop up throughout condensed matter physics.

The major topic of cond. mat.: phase transitions.

Quantum phase transitions (e.g. superconductivity, superfluidity) occur at zero or very low temperature characterised by abrupt change in ground state of system

→ quantum condensed matter  $\approx$  study of ground states of many-body systems.

The "Local Hamiltonian problem" asks whether or not a local Hamiltonian has a low energy ground state.

### Def (Local Hamiltonian Problem)

Input:  $k$ -local Hamiltonian  $H$  on  $n$  qudits with  $m$  local terms,

"Promise"  $\begin{cases} \text{where } \lambda_0(H) \leq \alpha \text{ or } \lambda_0(H) \geq \beta \\ \text{with } \beta - \alpha \geq \frac{1}{\text{poly}(n)} \end{cases}$

Output: YES if  $\lambda_0(H) \leq \alpha$   
NO if  $\lambda_0(H) \geq \beta$

Note that:

- Input is classical description of  $H$ , e.g. specify  $d^{2k}$  matrix elements for each of the  $m$  terms.
- Wlog  $m \leq \binom{n}{k} = O(n^k) = \text{poly}(n)$   
problem size (for fixed  $k, d$ ) scales as  $\text{poly}(n)$   
→ take  $n$  to measure problem size

- (Strictly speaking, number of bits of data required to specify input also depends on precision of matrix entries. All our results will hold if matrix entries are restricted to  $\text{poly}(n)$  digits of precision, so we ignore the input precision from now on.)
- Together with input, given promise that condition  $\lambda_0(H) \leq \alpha$  or  $\lambda_0(H) \geq \beta$ ,  $\beta - \alpha = \frac{1}{\text{poly}(n)}$  holds for the input  
 → only required to solve problem under assumption condition holds.
- Conversely, to prove hardness results must show condition does hold for any  $H$  we construct.

## Thm (Kitaev)

The Local Hamiltonian problem is QMA-hard

The first & most important result in Hamiltonian Complexity theory.

Has many interesting implications for QIT, CompSci, & especially physics (see later)

To prove QMA-hardness:

Transform any QMA prob.  $\rightarrow$  Local H. prob.

Recall Def. QMA:

$\exists$  poly-sized  $q$ . circuit  $U = U_T \cdot U_{T-1} \cdots U_2 \cdot U_1$  s.t.

YES:  $\exists |w\rangle : \Pr(U|w\rangle \text{ outputs "1"}) \geq 2/3$

NO:  $\forall |w\rangle : \Pr(U|w\rangle \text{ " "}) \leq 1/3$

Idea:

"Encode" verifier circuit  $\rightarrow$  local Hamiltonian s.t. ground state encodes output of circuit.

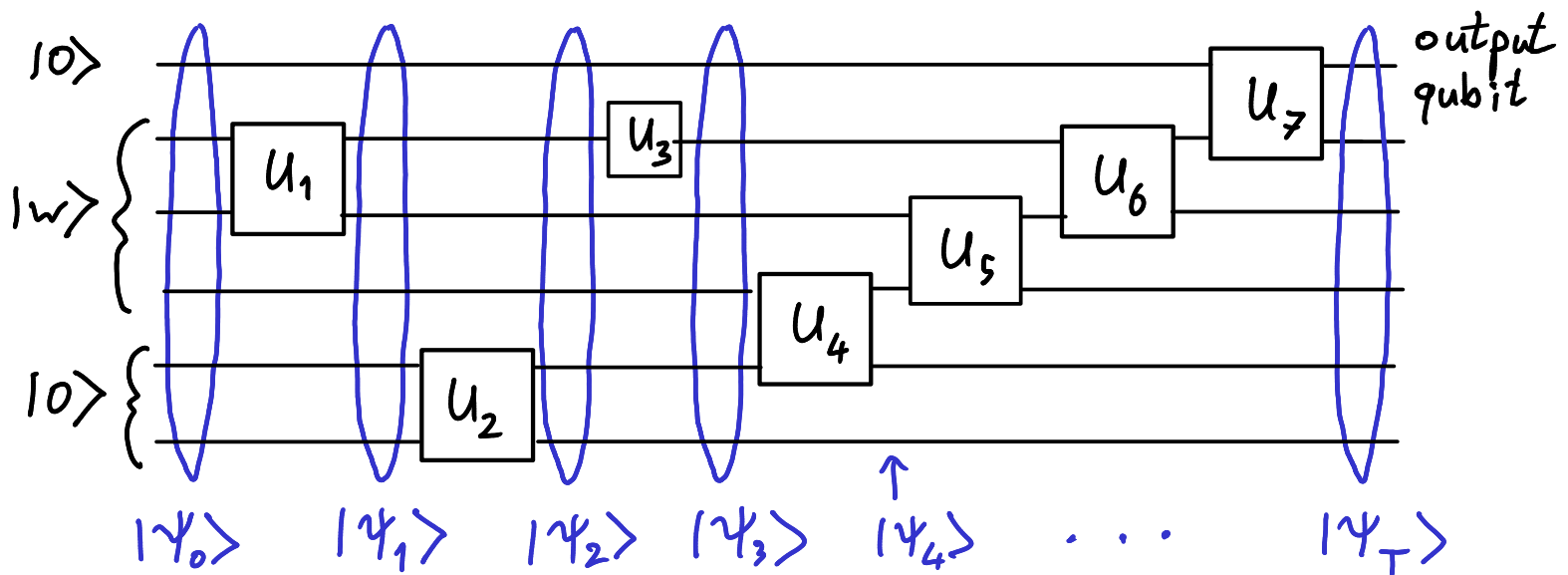
Add term that gives additional energy "penalty" if circuit outputs "0".

YES / NO  $\leftrightarrow$  low-energy / high energy g.s.

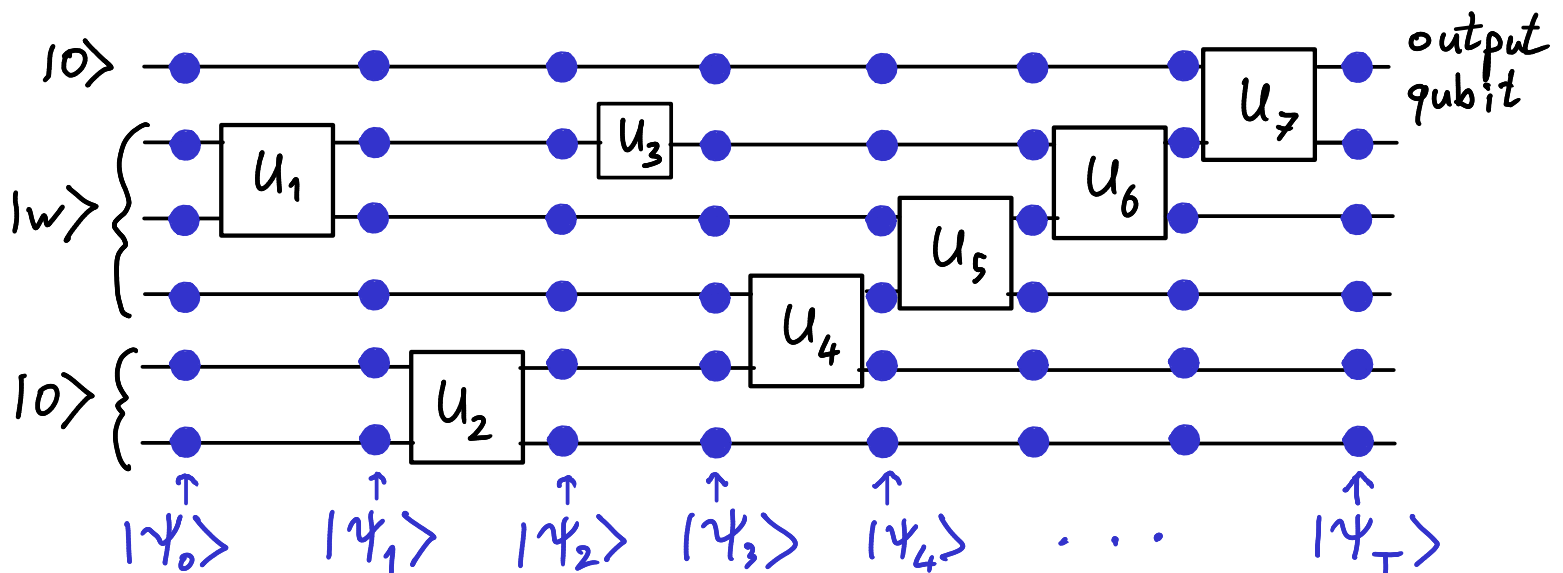
Before proving Thm, we first consider a naive approach that doesn't work, but teaches us something important about quantum many-body states and Hamiltonians.

Naive proof approach (doesn't work!)

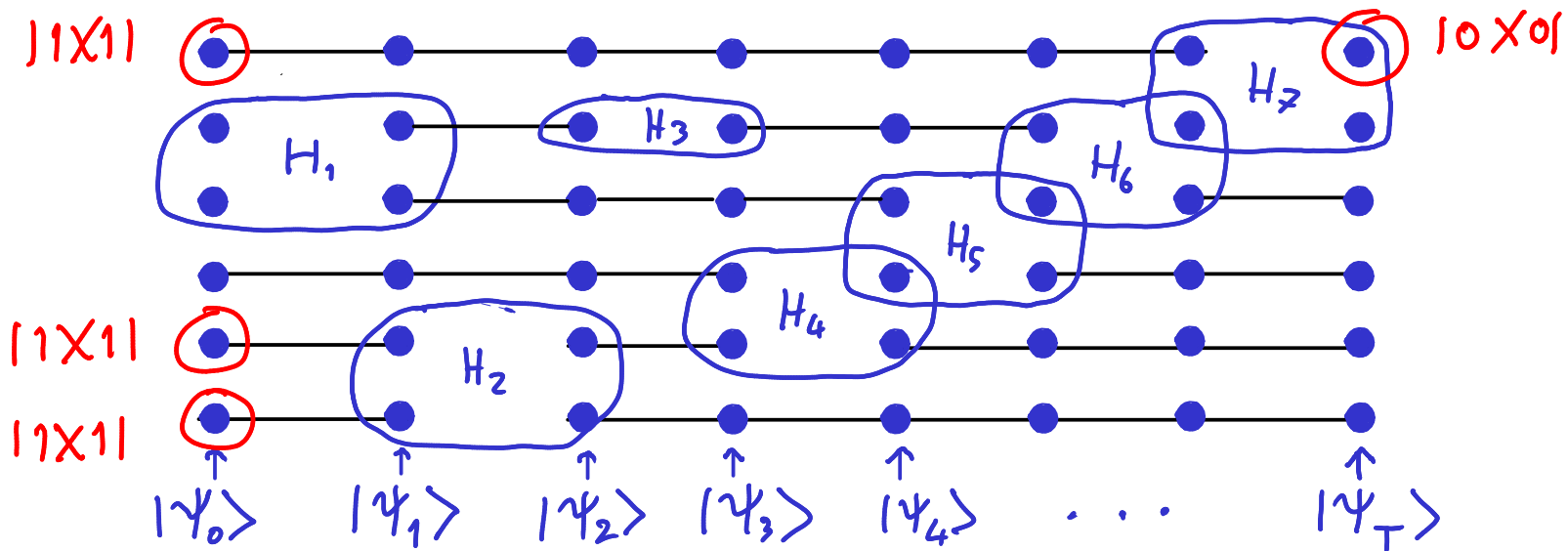
Verifier circuit:



Imagine laying out circuit on lattice of qubits:



Now put local Hamiltonians wherever there's a gate (also need 2-local Hamiltonian terms for identity gates):



Can easily force qubits representing initial state of output & ancilla qubits to be  $|0\rangle$  using 1-local term  $\Pi^{(1)} = |1X1\rangle$ ,

But leave initial witness qubits state unconstrained

Similarly, can give additional energy penalty if qubit representing final state of output qubit is "0" using 1-local term  $\Pi_{\text{out}}^{(0)} = |0X0\rangle$

If we could find Hamiltonian terms  $H_1, \dots, H_T$  s.t. ground states represent correct evolution of circuit, we'd be done:

YES instance:

$\exists$  witness  $|w\rangle$  s.t.  $\Pr(\text{output } |1\rangle) \geq \frac{2}{3}$

$\Rightarrow \exists$  state  $|\psi_0\rangle |\psi_1\rangle \dots |\psi_T\rangle$  representing correct evolution of circuit s.t. only picks up energy  $\leq \frac{1}{3}$  from  $\Pi_{\text{out}}^{(0)}$

$\rightarrow$  g.s. energy  $\leq \frac{1}{3}$

NO instance:

$\forall |w\rangle, \Pr(\text{output } |1\rangle) \leq \frac{1}{3}$

$\Rightarrow \forall |\psi\rangle = |\psi_0\rangle |\psi_1\rangle \dots |\psi_T\rangle$  representing correct evol<sup>n</sup> of circuit, pick up energy  $\geq \frac{2}{3}$  from  $\Pi_{\text{out}}^{(0)}$

$\forall |\psi\rangle$  not representing correct evol<sup>n</sup>, pick up large energy penalty from  $H_1, \dots, H_T$

$\rightarrow$  g.s. energy  $\geq \frac{2}{3}$

Exercise

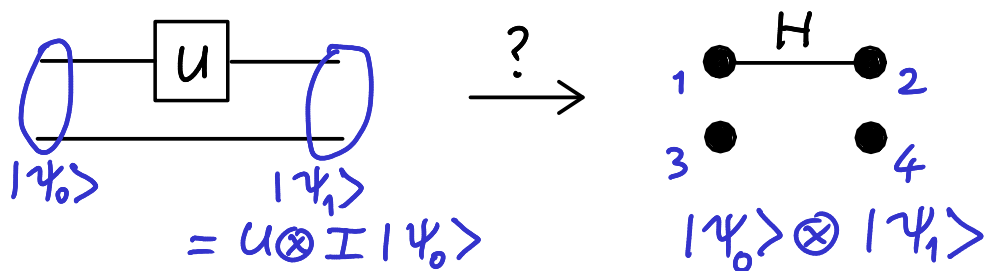
Make above argument rigorous.

If we could find Hamiltonian terms  $H_1, \dots, H_T$  s.t. ground states represent correct evolution of circuit, we'd be done.

### Exercise

Use this approach to prove NP-hardness of Local Hamiltonian problem for classical Hamiltonians (diagonal in computational basis).

Unfortunately, such Hamiltonian terms cannot exist in quantum case.



### Claim

$\exists U : \nexists H_{12}$  s.t. ground state subspace

$$\mathcal{L}_0(H_{12} \otimes \mathbb{1}_{34}) = \text{span} \{ |\psi\rangle_{13} \otimes (U \otimes \mathbb{1}) |\psi\rangle_{24} \}.$$

### Proof

wlog can take  $H \geq 0$ ,  $\lambda_{\min}(H) = 0$  (substitute  $H' = H + \lambda_{\min}(H) \cdot I$ , eigenvectors unchanged).

$$\rightarrow \mathcal{L}_0(H \otimes \mathbb{1}) = \ker(H \otimes \mathbb{1}).$$

Consider 2-qubit Hilbert space  $\mathcal{H} = \mathbb{C}^2 \otimes \mathbb{C}^2$  with trivial gate  $U = \mathbb{1}$ , so

$$\text{span} \{ |\psi\rangle_{13} \otimes (U \otimes \mathbb{1}) |\psi\rangle_{24} \} = \text{span} \{ |\psi\rangle_{13} |\psi\rangle_{24} \}.$$

Let  $|\phi\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$ . entangled state!  
(ebit)

$|\phi\rangle|\phi\rangle \in \ker(H \otimes \mathbb{1})$ , so

$$0 = \langle \phi |_{13} \langle \phi |_{24} (H_{12} \otimes \mathbb{1}_{34}) |\phi\rangle_{13} |\phi\rangle_{24}$$

$$= \text{tr} [ H_{12} \otimes \mathbb{1}_{34} \cdot |\phi\rangle_{13} |\phi\rangle_{24} \langle \phi |_{13} \langle \phi |_{24} ]$$

$$= \text{tr} [ H_{12} \cdot \text{tr}_3 (|\phi\rangle_{13} \langle \phi |_{13}) \otimes \text{tr}_4 (|\phi\rangle_{24} \langle \phi |_{24}) ]$$

$$= \text{tr} [ H_{12} \cdot \mathbb{1}_1 \otimes \mathbb{1}_2 ]$$

$$= \text{tr} H.$$

$H \geq 0$ ,  $\text{tr} H = 0 \Rightarrow H = 0$ .

$$\begin{aligned} \mathcal{L}_0(H \otimes \mathbb{1}) &= \ker(O \otimes \mathbb{1}) = \ker(O) \\ &= \mathcal{H} \neq \text{span} \{ |\psi\rangle |\psi\rangle \} \quad \square \end{aligned}$$

### Exercise

Show Proposition holds  $\forall U$ .